

Andersonian Deontic Logic, Propositional Quantification, and Mally

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Abstract We present a new axiomatization of the deontic fragment of Anderson's relevant deontic logic, give an Andersonian reduction of a relevant version of Mally's deontic logic previously discussed in this journal, study the effect of adding propositional quantification to Anderson's system, and discuss the meaning of Anderson's propositional constant in a wide range of Andersonian deontic systems.

1 Introduction

An Andersonian system of deontic logic is a system in which the deontic operator O ("it is obligatory that") is defined by $OA = e \Rightarrow A$, where e is a primitive propositional constant ("the good thing") and \Rightarrow is an implicational connective. The following systems are examples of Andersonian deontic systems.

1. The systems discussed in [1] and [2], in which \Rightarrow is strict implication, that is, $OA = \Box(e \rightarrow A)$, where \Box is the modal operator of necessity and \rightarrow is material implication.
2. Anderson's relevant deontic logic ([3], [4]), in which \Rightarrow is relevant implication.
3. The systems discussed in [8] and [14], in which \Rightarrow is strict relevant implication, that is, $OA = \Box(e \rightarrow A)$, where \Box is necessity and \rightarrow is relevant implication.

In this paper, we will present some new results on Andersonian deontic systems. We start with Anderson's relevant deontic logic. We give a new axiomatization of the deontic fragment of this system, show that the relevant version of Mally's deontic logic [13] presented in [11] is an extension of this fragment, and prove that e can be defined in terms of O as soon as propositional quantification is available. After

Received September 2, 2005; accepted January 2, 2006; printed November 14, 2006
2000 Mathematics Subject Classification: Primary, 03B47; Secondary, 03B45
Keywords: Anderson, deontic logic, relevance logic, propositional quantification, Mally
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this, we discuss some of the other Andersonian systems mentioned above. We will show that the addition of propositional quantification sheds light on the meaning of the constant e in these systems, too.

2 Anderson's Relevant Deontic Logic

Definition 2.1 (System \mathbf{R}) Relevant system \mathbf{R} has the following axioms and rules ([5], Ch. V).

(R1)	$A \rightarrow A$	(Self-implication)
(R2)	$(A \rightarrow B) \rightarrow ((C \rightarrow A) \rightarrow (C \rightarrow B))$	(Prefixing)
(R3)	$(A \rightarrow (B \rightarrow C)) \rightarrow (B \rightarrow (A \rightarrow C))$	(Permutation)
(R4)	$(A \rightarrow (A \rightarrow B)) \rightarrow (A \rightarrow B)$	(Contraction)
(R5)	$(A \& B) \rightarrow A, (A \& B) \rightarrow B$	(&Elimination)
(R6)	$((A \rightarrow B) \& (A \rightarrow C)) \rightarrow (A \rightarrow (B \& C))$	(&Introduction)
(R7)	$A \rightarrow (A \vee B), B \rightarrow (A \vee B)$	(\vee Introduction)
(R8)	$((A \rightarrow C) \& (B \rightarrow C)) \rightarrow ((A \vee B) \rightarrow C)$	(\vee Elimination)
(R9)	$(A \& (B \vee C)) \rightarrow ((A \& B) \vee C)$	(Distribution)
(R10)	$\neg\neg A \rightarrow A$	(Double Negation)
(R11)	$(A \rightarrow \neg B) \rightarrow (B \rightarrow \neg A)$	(Contraposition)
(\rightarrow E)	If A and $A \rightarrow B$ are theorems, B is a theorem	(Detachment)
(&I)	If A and B are theorems, $A \& B$ is a theorem	(Adjunction)

Definition: $A \leftrightarrow B = (A \rightarrow B) \& (B \rightarrow A)$.

Definition 2.2 (System \mathbf{R}_e) Anderson's relevant deontic logic \mathbf{R}_e is \mathbf{R} with a primitive propositional constant e and a unary operator O defined by $OA = e \rightarrow A$. Furthermore, the operator P is defined by $PA = \neg O\neg A$. e is read as "the good thing," O as "it is obligatory that," and P as "it is permitted that."¹

3 The Deontic Fragment of Anderson's Relevant Deontic Logic

To which "purely deontic" logic, stated in terms of O rather than e , does Anderson's proposal exactly give rise? This question has been answered by Goble [8].

Definition 3.1 (System $\mathbf{OR.1abc}$) Language: \mathbf{R} supplemented with a primitive propositional operator O . Axioms and rules: \mathbf{R} plus:

(OC)	$(OA \& OB) \rightarrow O(A \& B)$
(OK)	$O(A \rightarrow B) \rightarrow (OA \rightarrow OB)$
(ROa)	If $A \rightarrow B$ is a theorem, then so is $OA \rightarrow OB$
(a)	$O(OA \rightarrow A)$
(b)	$(A \rightarrow B) \rightarrow (OA \rightarrow OB)$
(c)	$A \rightarrow OPA$

Definition 3.2 (Deontic Fragment) The translation function h from the language of $\mathbf{OR.1abc}$ into the language of \mathbf{R}_e is defined as follows.

1. $h(A) = A$ if A is atomic
2. $h(\neg A) = \neg h(A)$
3. $h(A \& B) = h(A) \& h(B)$
4. $h(A \vee B) = h(A) \vee h(B)$
5. $h(A \rightarrow B) = h(A) \rightarrow h(B)$
6. $h(OA) = e \rightarrow h(A)$

The *deontic fragment* of \mathbf{R}_e (under h) is the set $\{A : \vdash_{\mathbf{R}_e} h(A)\}$.

Theorem 3.3 **OR.1abc** is an axiomatization of the deontic fragment of \mathbf{R}_e .

Proof By the Routley-Meyer semantics of \mathbf{R}_e and **OR.1abc**. See [8] for details.² \square

This result can be simplified, as we will now show.

Definition 3.4 (System \mathbf{R}_O) Language: \mathbf{R} supplemented with a primitive propositional operator O . Axioms and rules: \mathbf{R} plus:

- (a) $O(OA \rightarrow A)$
- (b) $(A \rightarrow B) \rightarrow (OA \rightarrow OB)$

Theorem 3.5 \mathbf{R}_O has the same theorems as **OR.1abc**.

Proof It is sufficient to prove that (ROa) is a derivable rule of \mathbf{R}_O and that (c), (OK), and (OC) are theorems of \mathbf{R}_O .

The proofs are as follows. We mention a few intermediate theorems for later reference. To avoid circularity, there are no forward references in this section.

(ROa) From (b) and detachment.

(Th1) $O(A \rightarrow B) \rightarrow (A \rightarrow OB)$.

- 1. $(A \rightarrow B) \rightarrow (A \rightarrow B)$ self-impl
- 2. $A \rightarrow ((A \rightarrow B) \rightarrow B)$ 1, permut
- 3. $A \rightarrow (O(A \rightarrow B) \rightarrow OB)$ 2, (b)
- 4. $O(A \rightarrow B) \rightarrow (A \rightarrow OB)$ 3, permut

(c) $A \rightarrow OPA$.

- 1. $O(O\neg A \rightarrow \neg A)$ (a)
- 2. $(O\neg A \rightarrow \neg A) \rightarrow (A \rightarrow PA)$ contrapos, def P
- 3. $O(A \rightarrow PA)$ 1, 2, (b)
- 4. $A \rightarrow OPA$ 3, (Th1)

(Th2) $(A \rightarrow OB) \rightarrow O(A \rightarrow B)$.

- 1. $(OB \rightarrow B) \rightarrow ((A \rightarrow OB) \rightarrow (A \rightarrow B))$ prefixing
- 2. $O(OB \rightarrow B) \rightarrow O((A \rightarrow OB) \rightarrow (A \rightarrow B))$ 1, (b)
- 3. $O((A \rightarrow OB) \rightarrow (A \rightarrow B))$ 2, (a)
- 4. $(A \rightarrow OB) \rightarrow O(A \rightarrow B)$ 3, (Th1)

(ROO) If OOA is a theorem, then OA is a theorem. Definition (for readability):

$D = (OA \rightarrow A)$.

- 1. OOA premise
- 2. $D \rightarrow (OOA \rightarrow OA)$ def D , (b)
- 3. $D \rightarrow OA$ 1, 2, permut
- 4. $D \rightarrow ((D \rightarrow OA) \rightarrow (D \rightarrow A))$ def D , pref
- 5. $D \rightarrow (D \rightarrow A)$ 3, 4, permut
- 6. $D \rightarrow A$ 5, contract
- 7. $OD \rightarrow OA$ 6, (b)
- 8. OD def D , (a)
- 9. OA 7, 8

(Th3) $OOA \rightarrow OA$.³

1. $OOA \rightarrow OOA$ self-impl
2. $O(OOA \rightarrow OA)$ 1, (Th2)
3. $OO(OOA \rightarrow A)$ 2, (Th2), (b)
4. $O(OOA \rightarrow A)$ 3, (ROO)
5. $OOA \rightarrow OA$ 4, (Th1)

(OK) $O(A \rightarrow B) \rightarrow (OA \rightarrow OB)$.

1. $O(A \rightarrow B) \rightarrow (A \rightarrow OB)$ (Th1)
2. $(A \rightarrow OB) \rightarrow (OA \rightarrow OOB)$ (b)
3. $(OA \rightarrow OOB) \rightarrow (OA \rightarrow OB)$ (Th3), pref
4. $O(A \rightarrow B) \rightarrow (OA \rightarrow OB)$ 1–3

(Th4) $(A \rightarrow B) \rightarrow (PA \rightarrow PB)$. From axiom (b) and (R10)–(R11).

(Th5) $POA \rightarrow A$. From theorem (c), axiom (b), and (R10)–(R11).

(OC) $(OA \& OB) \rightarrow O(A \& B)$.⁴

1. $P(OA \& OB) \rightarrow POA$ &Elim, (Th4)
2. $P(OA \& OB) \rightarrow A$ 1, (Th5)
3. $P(OA \& OB) \rightarrow POB$ &Elim, (Th4)
4. $P(OA \& OB) \rightarrow B$ 3, (Th5)
5. $P(OA \& OB) \rightarrow (A \& B)$ 2, 4, &Introd
6. $OP(OA \& OB) \rightarrow O(A \& B)$ 5, (b)
7. $(OA \& OB) \rightarrow O(A \& B)$ 6, (c)

This completes the proof. \square

Corollary 3.6 \mathbf{R}_O is an axiomatization of the deontic fragment of \mathbf{R}_e .

Interestingly enough, (OC) is not a theorem of *positive* \mathbf{R}_O , that is, positive \mathbf{R} with axioms (a) and (b) (proof: by MaGIC [18]).⁵ In other words, \mathbf{R}_O is not a conservative extension of positive \mathbf{R}_O . But we can state the following (by inspection of the proof in [8]).

Theorem 3.7 *Positive* \mathbf{R}_O plus (OC) is an axiomatization of the deontic fragment of positive \mathbf{R}_e .

In contrast with Goble's system **OR.1abc**, systems \mathbf{R}_O and positive \mathbf{R}_O are well axiomatized in the sense that the axioms are independent from each other (proof: by MaGIC [18]).

4 Mally's Deontic Logic

In an earlier paper [11], we presented the following relevant version of Ernst Mally's deontic logic, the first formal system of deontic logic ever put forward [13].

Definition 4.1 (System RD (Relevant Deontik)) Language: relevant system \mathbf{R} supplemented with a primitive unary operator O and a primitive propositional constant u ("the unconditionally obligatory"). Axioms and rules: \mathbf{R} plus:

- (I) $((A \rightarrow OB) \& (B \rightarrow C)) \rightarrow (A \rightarrow OC)$
- (II) $((A \rightarrow OB) \& (A \rightarrow OC)) \rightarrow (A \rightarrow O(B \& C))$
- (III) $(A \rightarrow OB) \leftrightarrow O(A \rightarrow B)$
- (IV) Ou
- (V) $\neg(u \rightarrow O\neg u)$

This system is the same as Mally's own system, except that Mally based his system on classical logic and accordingly accepted the "archetypical fallacy of relevance" $A \rightarrow (B \rightarrow A)$. Mally's system has the theorem $A \leftrightarrow OA$ (see [17]), but neither $A \rightarrow OA$ nor $OA \rightarrow A$ are theorems of **RD** (proof: by MaGIC [18]).

It can be shown that **RD** can be axiomatized more elegantly as follows (see [12] in combination with the derivation of theorem (OC) above).

- (I') $(A \rightarrow B) \rightarrow (OA \rightarrow OB)$
- (III') $O(OA \rightarrow A)$
- (IV') Ou
- (V') $\neg(u \rightarrow O\neg u)$

From the results we have presented above, it follows that **RD** is the deontic fragment of \mathbf{R}_e with an additional propositional constant u plus the axioms $e \rightarrow u$ and $\neg(u \rightarrow (e \rightarrow \neg u))$.

This result is useful because it makes it considerably easier to recognize some theorems of **RD** as such. For example, in [11] we failed to see that $OOA \rightarrow OA$ is a theorem, whereas we can now easily identify it as a theorem because $h(OOA \rightarrow OA)$, that is, $(e \rightarrow (e \rightarrow A)) \rightarrow (e \rightarrow A)$, is just an instance of Contraction.

5 Propositional Quantification (1)

What does the propositional constant e exactly mean? The following interpretations have been offered: "the good thing" or "good state of affairs" ([1], [2], [3], [4]), "what morality requires" [10], "optimality or admissibility" [7], "the content of an (unspecified) moral code" [19], "the law,"⁶ "it is not the case that all hell breaks loose" [20], "all normative demands are met" [16]. Most of these interpretations sound like poetry rather than logic. We shall show that once propositional quantifiers are added to \mathbf{R}_e , we can be more precise about the meaning of e .

Definition 5.1 (System $\mathbf{R}^{\forall p}$) Propositionally quantified relevant system $\mathbf{R}^{\forall p}$ has the following axioms and axiom clause in addition to those of **R** ([6], Ch. VI).

- (Q1) $\forall p(A \rightarrow B) \rightarrow (\forall pA \rightarrow \forall pB)$
- (Q2) $(\forall pA \ \& \ \forall pB) \rightarrow \forall p(A \ \& \ B)$
- (Q3) $\forall pA(p) \rightarrow A(B)$
- (Q4) $\forall p(A \rightarrow B) \rightarrow (A \rightarrow \forall pB)$ (p not free in A)
- (Q5) $\forall p(A \vee B) \rightarrow (A \vee \forall pB)$ (p not free in A)
- (Q*) If A is an axiom, then $\forall pA$ is an axiom.

Axioms Q1 and Q2 in conjunction with axiom clause Q* yield Generalization:

- (Gen) If A is a theorem, then $\forall pA$ is a theorem.

Definition 5.2 (System $\mathbf{R}_e^{\forall p}$) $\mathbf{R}_e^{\forall p}$ is $\mathbf{R}^{\forall p}$ with primitive propositional constant e and propositional operators O and P defined as in \mathbf{R}_e .

Theorem 5.3 $\mathbf{R}_e^{\forall p}$ has the following theorem: $e \leftrightarrow \forall p(Op \rightarrow p)$.

Proof

1.	$(e \rightarrow A) \rightarrow (e \rightarrow A)$	self-impl	
2.	$e \rightarrow ((e \rightarrow A) \rightarrow A)$	1, permut	
3.	$e \rightarrow \forall p((e \rightarrow p) \rightarrow p)$	2, (Gen), (Q4)	
4.	$e \rightarrow \forall p(Op \rightarrow p)$	3, def O	
5.	$\forall p(Op \rightarrow p) \rightarrow \forall p((e \rightarrow p) \rightarrow p)$	def O	
6.	$\forall p((e \rightarrow p) \rightarrow p) \rightarrow ((e \rightarrow e) \rightarrow e)$	(Q3)	
7.	$(e \rightarrow e)$	self-impl	
8.	$((e \rightarrow e) \rightarrow e) \rightarrow e$	7, self-impl, permut	
9.	$\forall p(Op \rightarrow p) \rightarrow e$	5, 6, 8	
10.	$e \leftrightarrow \forall p(Op \rightarrow p)$	4, 9, adj	□

Thus e says that all obligations are fulfilled (all normative demands are met). This happens to agree with McNamara's unmotivated informal reading of e [16].

6 Inverse Andersonian Reduction (1)

Theorem 5.3 suggests the following question: is it possible to define e in terms of O and carry out the inverse of the Andersonian reduction? We shall show that this is indeed possible.

Definition 6.1 (System $\mathbf{R}_O^{\forall p}$) System $\mathbf{R}_O^{\forall p}$ is $\mathbf{R}^{\forall p}$ supplemented with a primitive operator O , a propositional constant e defined by $e = \forall p(Op \rightarrow p)$, and axioms (a) and (b) of \mathbf{R}_O .

Theorem 6.2 $\mathbf{R}_O^{\forall p}$ and $\mathbf{R}_e^{\forall p}$ have the same theorems.

Proof First, all theorems of $\mathbf{R}_O^{\forall p}$ are theorems of $\mathbf{R}_e^{\forall p}$. All cases are easy except perhaps $e \leftrightarrow \forall p(Op \rightarrow p)$ (which is a theorem of $\mathbf{R}_O^{\forall p}$ because of the definition of e in $\mathbf{R}_O^{\forall p}$), but this case has already been discussed (Theorem 5.3).

Second, all theorems of $\mathbf{R}_e^{\forall p}$ are theorems of $\mathbf{R}_O^{\forall p}$. It is sufficient to prove that $OA \leftrightarrow (e \rightarrow A)$ is a theorem of $\mathbf{R}_O^{\forall p}$. (This formula is a theorem of $\mathbf{R}_e^{\forall p}$ because of the definition of O in $\mathbf{R}_e^{\forall p}$.)

Let us first derive the deontic Barcan formula:

$$(OBF) \quad \forall pOA(p) \rightarrow O\forall pA(p).$$

The following derivation is similar to the proof of (OC) above. It is also similar to a well-known proof of the Barcan formula in individually quantified **S5** ([9], p. 247).

1.	$P\forall pOA(p) \rightarrow POA(B)$	Q3, (Th4)
2.	$P\forall pOA(p) \rightarrow A(B)$	1, (Th5)
3.	$\forall q(P\forall pOA(p) \rightarrow A(q))$	2, (Gen)
4.	$P\forall pOA(p) \rightarrow \forall pA(p)$	3, (Q4)
5.	$OP\forall pOA(p) \rightarrow O\forall pA(p)$	4, (b)
6.	$\forall pOA(p) \rightarrow O\forall pA(p)$	5, (c)

We can now derive $OA \leftrightarrow (e \rightarrow A)$.

- | | | | |
|----|--|----------------|---|
| 1. | $e \rightarrow (OA \rightarrow A)$ | def e , (Q3) | |
| 2. | $OA \rightarrow (e \rightarrow A)$ | 1, permut | |
| 3. | $O(e \rightarrow A) \rightarrow (Oe \rightarrow OA)$ | (OK) | |
| 4. | $Oe \rightarrow ((e \rightarrow A) \rightarrow OA)$ | 3, permut | |
| 5. | $\forall p O(Op \rightarrow p)$ | (a), (Q*) | |
| 6. | $O\forall p(Op \rightarrow p)$ | 5, (OBF) | |
| 7. | Oe | 6, def e | |
| 8. | $(e \rightarrow A) \rightarrow OA$ | 4, 7 | |
| 9. | $OA \leftrightarrow (e \rightarrow A)$ | 2, 8, adj | □ |

We might call this “a reduction of alethic logic to deontic logic.”

7 Weaker Andersonian Systems

Some authors ([8], [15], [14]) have objected to axiom (b) and theorems (c), (Th1), and (Th2) of \mathbf{R}_O . An alternative approach in the Andersonian tradition is as follows ([8], [14]): start from \mathbf{R} , add the modal operator of necessity \Box and some axioms and rules for \Box , add the constant e , and define O by $OA = \Box(e \rightarrow A)$. In the resulting deontic systems, (b), (c), (Th1), and (Th2) are, in general, not derivable. In the following, we shall study to which extent the results obtained for Anderson’s relevant deontic logic are valid for these weaker systems.

8 Propositional Quantification (2)

Andersonian systems \mathbf{RT}_{Vmo} [14] and $\mathbf{RT.a}$ [8] have the following axiom and rule (among others) in addition to those of R , along with the just-mentioned definition of O :

- (□T) $\Box A \rightarrow A$
(Nec) If A is a theorem, $\Box A$ is a theorem.

Principles (□T) and (Nec) suffice to derive $e \leftrightarrow \forall p(Op \rightarrow p)$ in the propositionally quantified versions of these Andersonian systems. The proof is as follows.

- | | | |
|-----|--|----------------------|
| 1. | $\Box(e \rightarrow A) \rightarrow (e \rightarrow A)$ | (□T) |
| 2. | $e \rightarrow (\Box(e \rightarrow A) \rightarrow A)$ | 1, permut |
| 3. | $e \rightarrow \forall p(\Box(e \rightarrow p) \rightarrow p)$ | 2, (Gen), (Q4) |
| 4. | $e \rightarrow \forall p(Op \rightarrow p)$ | 3, def O |
| 5. | $\forall p(Op \rightarrow p) \rightarrow \forall p(\Box(e \rightarrow p) \rightarrow p)$ | def O |
| 6. | $\forall p(\Box(e \rightarrow p) \rightarrow p) \rightarrow (\Box(e \rightarrow e) \rightarrow e)$ | (Q3) |
| 7. | $\Box(e \rightarrow e)$ | self-impl, (Nec) |
| 8. | $(\Box(e \rightarrow e) \rightarrow e) \rightarrow e$ | 7, self-impl, permut |
| 9. | $\forall p(Op \rightarrow p) \rightarrow e$ | 5, 6, 8 |
| 10. | $e \leftrightarrow \forall p(Op \rightarrow p)$ | 4, 9, adj |

Most mixed alethic-deontic systems (both relevant and classical) discussed in the literature ([7], [8]) have (□T) and are closed under (Nec), with the result that $e \leftrightarrow \forall p(Op \rightarrow p)$ is a theorem of the propositionally quantified versions of these systems. It is appropriate to read e as “all obligations are fulfilled” in this whole range of systems. In weaker systems, however, one may have a greater freedom of interpretation.

9 Inverse Andersonian Reduction (2)

As before, we may ask whether we can carry out the inverse of the Andersonian reduction in the deontic fragments of these weaker Andersonian systems and define e in terms of O and the other connectives. The answer is affirmative, provided that one is considering sufficiently strong mixed alethic-deontic propositionally quantified systems.

Definition 9.1 (System $\mathbf{RS4}$) Relevant alethic modal system $\mathbf{RS4}$ has the following axioms and rules in addition to those of \mathbf{R} .

- ($\square\mathbf{K}$) $\square(A \rightarrow B) \rightarrow (\square A \rightarrow \square B)$
- ($\square\mathbf{C}$) $(\square A \ \& \ \square B) \rightarrow \square(A \ \& \ B)$
- ($\square\mathbf{T}$) $\square A \rightarrow A$
- ($\square\mathbf{4}$) $\square A \rightarrow \square\square A$
- (Nec) If A is a theorem, then $\square A$ is a theorem.

Definition 9.2 (System $\mathbf{RS4}_e$) System $\mathbf{RS4}_e$ has the same axioms and rules as $\mathbf{RS4}$ but it also contains a primitive propositional constant e and a propositional operator O defined by $OA = \square(e \rightarrow A)$.

Definition 9.3 (System $\mathbf{RS4}_O$) Mixed alethic-deontic system $\mathbf{RS4}_O$ is $\mathbf{RS4}$ supplemented with a primitive operator O and the following axioms in addition to the axioms and rules of $\mathbf{RS4}$.

- (\mathbf{OK}) $\square(A \rightarrow B) \rightarrow (OA \rightarrow OB)$
- (\mathbf{OC}) $(OA \ \& \ OB) \rightarrow O(A \ \& \ B)$
- (\mathbf{OT}) $O(OA \rightarrow A)$
- ($\mathbf{O4}$) $OA \rightarrow \square OA$

The notion “deontic fragment” is defined as above (Definition 3.2), except that the clause for O is changed to $h(OA) = \square(e \rightarrow h(A))$ and the clause $h(\square A) = \square(h(A))$ is added.

Theorem 9.4 $\mathbf{RS4}_O$ is an axiomatization of the deontic fragment of $\mathbf{RS4}_e$.

Proof By the Routley-Meyer semantics of $\mathbf{RS4}_O$ and $\mathbf{RS4}_e$. See [8] for details. \square

Definition 9.5 (System $\mathbf{RS4}_e^{\forall p}$) Propositionally quantified relevant alethic modal system $\mathbf{RS4}_e^{\forall p}$ has the following axiom in addition to the axioms and rules of $\mathbf{R}^{\forall p}$ and $\mathbf{RS4}_e$.

- ($\square\mathbf{BF}$) $\forall p \square A \rightarrow \square \forall p A$

The operator O is defined as in $\mathbf{RS4}_e$.

Definition 9.6 (System $\mathbf{RS4}_O^{\forall p}$) Propositionally quantified relevant mixed alethic-deontic system $\mathbf{RS4}_O^{\forall p}$ has the following axiom in addition to the axioms and rules of $\mathbf{R}^{\forall p}$ and $\mathbf{RS4}_O$.

- (\mathbf{OBF}) $\forall p OA \rightarrow O \forall p A$

$\mathbf{RS4}_O^{\forall p}$ also contains a constant expression e defined by $e = \forall p(O p \rightarrow p)$.

Theorem 9.7 $\mathbf{RS4}_O^{\forall p}$ and $\mathbf{RS4}_e^{\forall p}$ have the same theorems.

Proof First, all theorems of $\mathbf{RS4}_O^{\forall p}$ are theorems of $\mathbf{RS4}_e^{\forall p}$. All cases are easy except perhaps $e \leftrightarrow \forall p(O p \rightarrow p)$, which has already been discussed (Section 8).

Second, all theorems of $\mathbf{RS4}_e^{\forall p}$ are theorems of $\mathbf{RS4}_O^{\forall p}$. It is sufficient to prove that $OA \leftrightarrow \Box(e \rightarrow A)$ is a theorem of $\mathbf{RS4}_O^{\forall p}$.

- | | | | |
|-----|---|-------------------|---|
| 1. | $e \rightarrow (OA \rightarrow A)$ | def e , (Q3) | |
| 2. | $OA \rightarrow (e \rightarrow A)$ | 1, permut | |
| 3. | $\Box(OA \rightarrow (e \rightarrow A))$ | 2, (Nec) | |
| 4. | $\Box OA \rightarrow \Box(e \rightarrow A)$ | 3, (\Box K) | |
| 5. | $OA \rightarrow \Box(e \rightarrow A)$ | 4, (O4) | |
| 6. | $Oe \rightarrow (\Box(e \rightarrow A) \rightarrow OA)$ | (OK), permut | |
| 7. | $\forall p O(Op \rightarrow p)$ | (OT), (Q*) | |
| 8. | Oe | 7, (OBF), def e | |
| 9. | $\Box(e \rightarrow A) \rightarrow OA$ | 6, 8 | |
| 10. | $OA \leftrightarrow \Box(e \rightarrow A)$ | 5, 9, adj | □ |

Axiom (OC) of $\mathbf{RS4}_O^{\forall p}$ is not used in this derivation of $OA \leftrightarrow \Box(e \rightarrow A)$. On the other hand, (OC) is a theorem in the presence of $OA \leftrightarrow \Box(e \rightarrow A)$:

- | | | |
|----|---|--|
| 1. | $(OA \& OB) \rightarrow (\Box(e \rightarrow A) \& \Box(e \rightarrow B))$ | $OA \leftrightarrow \Box(e \rightarrow A)$ |
| 2. | $(\Box(e \rightarrow A) \& \Box(e \rightarrow B)) \rightarrow \Box(e \rightarrow (A \& B))$ | \Box C, &Intr, \Box K |
| 3. | $\Box(e \rightarrow (A \& B)) \rightarrow O(A \& B)$ | $OA \leftrightarrow \Box(e \rightarrow A)$ |
| 4. | $(OA \& OB) \rightarrow O(A \& B)$ | 1–3 |

This means that axiom (OC) of $\mathbf{RS4}_O^{\forall p}$ is redundant, even though it is not redundant in $\mathbf{RS4}_O$ (proof: by MaGIC [18]).

Theorem 9.7 also holds for stronger mixed alethic-deontic systems, such as systems based on classical **S4**.

10 Conclusion

We have demonstrated the following four facts. First, the deontic fragment of Anderson's relevant deontic logic has a very short and simple axiomatization. Second, Mally's deontic system [13] as reformulated in [11] is the deontic fragment of an extension of Anderson's relevant deontic logic. Third, the propositionally quantified versions of Anderson's own relevant deontic system and most other Andersonian deontic systems proposed in the literature have the theorem $e \leftrightarrow \forall p(O p \rightarrow p)$. In all these systems it is therefore appropriate to read e as "all obligations are fulfilled." Fourth, in some sufficiently strong deontic systems it is possible to define e by $e = \forall p(O p \rightarrow p)$ and then prove an Andersonian reduction principle of the form $OA \leftrightarrow (e \rightarrow A)$ or $OA \leftrightarrow \Box(e \rightarrow A)$, where \rightarrow is relevant or material implication. This provides a second justification for reading e as "all obligations are fulfilled."

Notes

1. Anderson also considered an "axiom of avoidance" $\neg(e \rightarrow \neg e)$ or $OA \rightarrow PA$ (see [8], [15], and [14]), which we will ignore.
2. The author would like to thank Lou Goble for making [8] available to him and discussing it with him.

3. An alternative derivation of (Th3) is to be found in [12].
4. The proof of (OC) presented here is due to John Slaney. We thank him for his generous permission to reproduce it here, and Bob Meyer for asking him to produce it. The proof is similar to the proof of the deontic Barcan formula (OBF) in Section 6, which is in turn similar to the “classical” proof of the Barcan formula in individually quantified **S5** ([9], p. 247).
5. This result was first obtained by Meyer, shared in private communication.
6. From private communication with Bob Meyer.

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Acknowledgments

The author is very grateful to Lou Goble, Bob Meyer, and John Slaney for discussing this paper with him and for giving him permission to include some of their own unpublished results in it.

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