

# The Logic of Common Ignorance

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# Introduction

**Quote** “Knowledge is a big subject. Ignorance is bigger. . . and it is more interesting.”<sup>1</sup>

**Claim** Ignorance has some surprising properties.

**Example** Common ignorance.

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<sup>1</sup>Stuart Firestein, Interview about S. Firestein, *Ignorance: How It Drives Science*, OUP 2012.


## Question

- ▶ “Obama calls Trump ignorant about foreign affairs” (Google, August 16, 2016, 8 results).
- ▶ “Trump calls Obama ignorant about foreign affairs” (Google, August 16, 2016, about 135 results).
- ▶ Suppose that at least one of them were right. (Of course, both could be right.)
- ▶ Would this give the group of all humans *common ignorance* about foreign affairs?

## Knowing that

To answer this question, we extend the (propositional) logic of individual, shared and common knowledge that  $A$ ,  $TEC_{(m)}$ , with a few uncontroversial definitions.  $TEC_{(m)}$  applies to a group having members  $1, \dots, m$ .  $TEC_{(m)}$  is well-known and is axiomatized as follows.<sup>2</sup>

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<sup>2</sup>J.-J. Ch. Meyer and W. van der Hoek, *Epistemic Logic for Computer Science and Artificial Intelligence* (Cambridge: Cambridge University Press, 1995), Ch. 2.1. 

# Symbols

- ▶ Individual knowledge that  $A$ :  $K_i A$ , where  $1 \leq i \leq m$ .  $K_i A$  is read as “ $i$  individually knows that  $A$ ” or as “ $i$  has individual knowledge that  $A$ .”
- ▶ Shared knowledge that  $A$ :  $EA$ .  $EA$  is read as “everyone knows that  $A$ ” or as “the group has shared knowledge that  $A$ .”
- ▶ Common knowledge that  $A$ :  $CA$ .  $CA$  is read as “it is commonly known that  $A$ ” or as “the group has common knowledge that  $A$ .”

# Axioms and derivation rules

A1 All instances of propositional tautologies.

A2  $\mathbf{K}_i(A \rightarrow B) \rightarrow (\mathbf{K}_iA \rightarrow \mathbf{K}_iB)$ .

A3  $\mathbf{K}_iA \rightarrow A$ .

A4  $\mathbf{E}A \leftrightarrow \bigwedge_{i=1}^m \mathbf{K}_iA$ .

A5  $\mathbf{C}A \rightarrow A$ .

A6  $\mathbf{C}A \rightarrow \mathbf{E}CA$ .

A7  $\mathbf{C}(A \rightarrow B) \rightarrow (\mathbf{C}A \rightarrow \mathbf{C}B)$ .

A8  $\mathbf{C}(A \rightarrow \mathbf{E}A) \rightarrow (A \rightarrow \mathbf{C}A)$ .

R1 From  $A$  and  $A \rightarrow B$  infer  $B$ .

R2 From  $A$  infer  $\mathbf{K}_iA$ .

R3 From  $A$  infer  $\mathbf{C}A$ .

# Theorems

- 1.1  $\mathbf{CA} \rightarrow \mathbf{EA}$  (common knowledge that  $A$  implies shared knowledge that  $A$ ).
- 1.2  $\mathbf{EA} \rightarrow \mathbf{K}_i A$  (shared knowledge that  $A$  implies individual knowledge that  $A$ ).
- 1.3  $\mathbf{CA} \rightarrow \mathbf{K}_i A$  (common knowledge that  $A$  implies individual knowledge that  $A$ ).
- †1.4  $\mathbf{K}_i A \rightarrow \mathbf{CA}$  (individual knowledge that  $A$  implies common knowledge that  $A$ ) is *invalid* [proof: by the semantics].

Intuitively,  $\mathbf{CA} = \bigwedge_{i \geq 0} \mathbf{E}^i A$  (common knowledge that  $A$  is the conjunction of  $A$ , shared knowledge that  $A$ , shared knowledge that the group has shared knowledge that  $A$ , and so on).

# Knowledge whether/about

Symbols:<sup>3</sup>

- ▶ Individual knowledge about  $A$ :  $\Delta_i A = K_i A \vee K_i \neg A$ .  $\Delta_i A$  is read as “ $i$  individually knows whether  $A$ ” or as “ $i$  has individual knowledge about  $A$ .”
- ▶ Common knowledge about  $A$ :  $C_{\Delta} A = C A \vee C \neg A$ .  $C_{\Delta} A$  is read as “the group has common knowledge about  $A$ .”

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<sup>3</sup>See J. Fan, Y. Wang and H. van Ditmarsch, “Contingency and Knowing Whether,” *The Review of Symbolic Logic*, 8:75–107, 2015.



# Theorems

2.1  $\mathbf{C}_\Delta A \rightarrow \Delta_i A$  [ $(\mathbf{C}A \vee \mathbf{C}\neg A) \rightarrow (\mathbf{K}_i A \vee \mathbf{K}_i \neg A)$ ] (common knowledge about  $A$  implies individual knowledge about  $A$ )  
[from  $\mathbf{C}A \rightarrow \mathbf{K}_i A$  (1.3) by propositional calculus].

†2.2  $\Delta_i A \rightarrow \mathbf{C}_\Delta A$  (individual knowledge about  $A$  implies common knowledge about  $A$ ) is *invalid* [proof: by the semantics].

# Ignorance whether/about

Symbols:<sup>4</sup>

- ▶ Individual ignorance about  $A$ :

$\nabla_i A = \neg \Delta_i A = \neg K_i A \wedge \neg K_i \neg A$  (individual ignorance about  $A$  is the negation of individual knowledge about  $A$ ).  $\nabla_i A$  is read as “ $i$  does not individually know whether  $A$ ”, as “ $i$  individually ignores whether  $A$ ” or as “ $i$  has individual ignorance about  $A$ .”

- ▶ Common ignorance about  $A$ :

$C_{\nabla} A = \neg C_{\Delta} A = \neg C A \wedge \neg C \neg A$  (common ignorance about  $A$  is the negation of common knowledge about  $A$ ).  $C_{\nabla} A$  is read as “the group has common ignorance about  $A$ .”

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<sup>4</sup>See Fan, Wang and Van Ditmarsch, “Contingency and Knowing Whether,”  
op. cit.

# Theorems

3.1  $\nabla_i A \rightarrow \mathbf{C}_{\nabla} A$  [ $\neg \Delta_i A \rightarrow \neg \mathbf{C}_{\Delta} A$ ] (individual ignorance about  $A$  implies common ignorance about  $A$ ) [from  $\mathbf{C}_{\Delta} A \rightarrow \Delta_i A$  (2.1) by contraposition].

†3.2  $\mathbf{C}_{\nabla} A \rightarrow \nabla_i A$  (common ignorance about  $A$  implies individual ignorance about  $A$ ) is *invalid* [proof: by the semantics].

Individual ignorance about  $A$  is therefore stronger than common ignorance about  $A$ . If agents have individual ignorance about  $A$ , all groups to which they belong have common ignorance about  $A$ .

## Answer to question

- ▶ Obama and Trump called each other ignorant about foreign affairs.
- ▶ Suppose that at least one of them were right.
- ▶ Question: would this give the group of all humans *common ignorance* about foreign affairs?
- ▶ Answer: yes, it would, by theorem  $\nabla_i A \rightarrow \mathbf{C}_{\nabla} A$  (3.1).

## Common ignorance about common ignorance

- ▶  $S5EC_{(m)}$  is  $TEC_{(m)}$  plus  $\neg K_i A \rightarrow K_i \neg K_i A$  (“ $i$  does not know that  $A$ ” implies “ $i$  knows that  $i$  does not know that  $A$ ”).
- ▶  $S5EC_{(m)}$  has the following theorem.<sup>5</sup>
  - 4.1  $\neg C_{\nabla} C_{\nabla} A$  (there is no common ignorance about common ignorance about  $A$ ).
- ▶  $TEC_{(m)}$  does not have this theorem, as the semantics shows.
- ▶ The Obama/Trump case seems to show that 4.1 is false.
- ▶ We do have common ignorance about our common ignorance about foreign affairs.
- ▶  $TEC_{(m)}$  is therefore preferable to  $S5EC_{(m)}$ .

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<sup>5</sup>H. Montgomery and R. Routley, “Contingency and Non-Contingency Bases for Normal Modal Logics,” *Logique et Analyse*, 9:318–328, 1966.